

## podft: On Accelerating Dynamic Taint Analysis with Precise Path Optimization

Zhiyou Tian Xidian University 21151213645@stu.xidian.edu.cn Cong Sun Xidian University suncong@xidian.edu.cn

#### Dongrui Zeng

Palo Alto Networks dzeng@paloaltonetworks.com Gang Tan Pennsylvania State University gtan@psu.edu

BAR 2023



- Dynamic taint analysis (DTA)
  - What is it?
  - Useful for security





- Binary-level dynamic data-flow tracking (DFT)
  - Dynamic binary instrumentation (DBI)
  - Virtual machine manager (VMM)
  - Emulator





- DBI-based DTA
  - Focus on explicit flows
  - Hold the tainting states within tagging memory

• Challenge of DTA —— significant performance penalty





历安冠子科技大学 XIDIAN UNIVERSITY

- Existing works
  - Lift (*MICOR 2006*)
    - static fast path
  - Libdft (VEE 2012)
    - on Pin
    - DBI inline routines
  - TaintRabbit (ASIA CCS 2020)
    - on DynamoRIO
    - dynamic fast path
  - SELECTIVETAINT (USENIX 2021)
    - static binary rewriting bloat the attack surface
    - value-set analysis cannot work on library code

Our work — — podft defines and enforces various fast paths

## podft advantages

- more efficient
- not bloat the attack surface
- consider library code
- flexible scalability



#### Design of podft

### podft overview

- BPA-based CFG Construction
- VSA-based tainted inst identification
- Tracking policy construction
- PDG-based function abstract(from SDFT)
- Pin-based Tracker

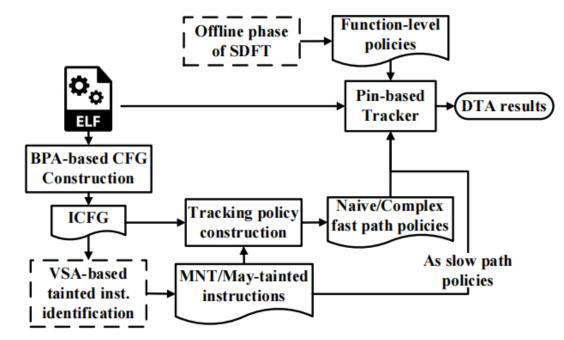


Fig1. Framework of podft (dashed block = usage of existing tools)

#### Next —— give a toy example to demonstrate

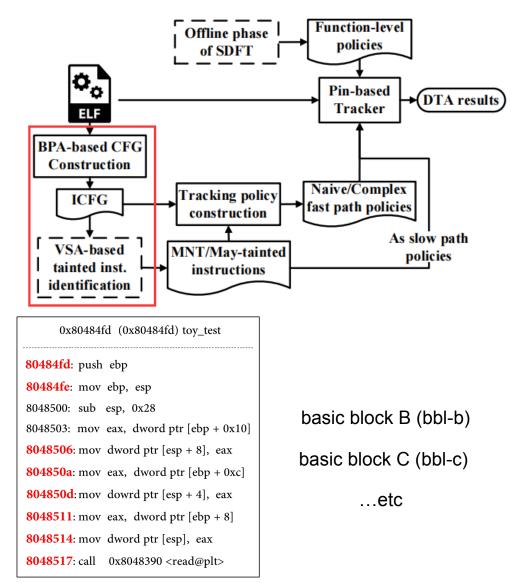




- BPA-based CFG Construction
- VSA-based tainted inst identification

```
void toy_test (int fd, char *buf, int size){
    int read_len = read(fd, buf, size); // taint source
    if (read_len > 0){
        printf ("read data: %s\n", buf);
        for (int i = 0; i<2; i++){
            buf[i] = i;
            write (fd, buf[i], 1); // taint sink
        }
    }
}
int main(int argc, char *argv[]){
    char buffer [64] = {0};
    int fd = open(argv[1], O_RDWR);
    toy_test (fd, buffer, 64);
    return 0;
}</pre>
```

Fig.2 toy example





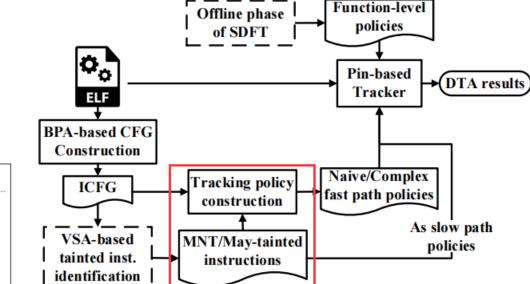
#### Tracking policy construction

- naive fast path → main(bbl-c)
- complex fast path → toy\_test+0x44(bbl-b)
- slow path →toy\_test(bbl-a)
- function fast path → printf etc..

0x80484fd (0x80484fd) toy_test
80484fd: push ebp
80484fe: mov ebp, esp
8048500: sub esp, 0x28
8048503: mov eax, dword ptr [ebp + 0x10]
8048506: mov dword ptr [esp + 8], eax
<b>804850a</b> : mov eax, dword ptr [ebp + 0xc]
804850d: mov dowrd ptr [esp + 4], eax
8048511: mov eax, dword ptr [ebp + 8]
8048514: mov dword ptr [esp], eax
8048517: call 0x8048390 <read@plt></read@plt>

0x8048541 (0x80484fd) toy\_test+0x44 **8048541**: mov edx, dword ptr [ebp - 0x10] **8048544**: mov eax, dword ptr [ebp + 0xc] 8048547: add edx, eax **8048549**: mov eax, dword ptr [ebp – 0x10] 804854c: mov byte ptr [edx], al **804854e**: mov edx, dword ptr [ebp – 0x10] **8048551**: mov eax, dword ptr [ebp + 0xc] 8048554: add eax, edx 8048556: movzx eax, byte ptr [eax] 8048559: movsx eax, al **804855c**: mov dword ptr [esp + 8], 1 8048564: mov dword ptr [esp + 4], eax **8048568**: mov eax, dword ptr [ebp + 8] 804856b: mov dword ptr [esp], eax **804856e**: call 0x80483f0 <write@plt>

0x804857f (0x804857f) main 804857f: push ebp 8048580: mov ebp, esp 8048582: push edi 8048583: push esi 8048584: push ebx, 8048585: and esp, 0xfffffff0 8048588: sub esp, 0x60 804858b: mov eax, dword ptr [ebp + 0xc]804858e: mov dword ptr [esp + 0xc], eax 8048592: mov eax, dword ptr gs : [0x14] 8048598: mov dword ptr [esp + 0x5c], eax 804859c: xor eax, eax 804859e: lea ebx, [esp + 0x1c]80485a2: mov eax, 0 80485a7: mov edx, 0x10 80485ac: mov edi, ebx 80485ae: mov ecx, edx 80485b0: rep stosd dword ptr es : [edi], eax



Why?

basic block A (bbl-a)

basic block B (bbl-b)

basic block C (bbl-c)





#### • Tracking policy construction

- naive fast path → main(bbl-c)
  - Not contain potentially tainted instructions

# because

0x804857f (0x804857f) main 804857f: push ebp 8048580: mov ebp, esp 8048582: push edi 8048583: push esi 8048584: push ebx, 8048585: and esp, 0xfffffff0 8048588: sub esp, 0x60 804858b: mov eax, dword ptr [ebp + 0xc]804858e: mov dword ptr [esp + 0xc], eax 8048592: mov eax, dword ptr gs : [0x14] 8048598: mov dword ptr [esp + 0x5c], eax 804859c: xor eax, eax 804859e: lea ebx, [esp + 0x1c]80485a2: mov eax, 0 80485a7: mov edx, 0x10 80485ac: mov edi, ebx 80485ae: mov ecx, edx 80485b0: rep stosd dword ptr es : [edi], eax

basic block C (bbl-c)



#### Tracking policy construction

- complex fast path → toy\_test+0x44(bbl-b)
  - Contain potentially tainted instructions
  - Hot BBL (be executed multiple times)
  - TaintedMem(bbl)  $\cap$  MergedDep(bbl) = $\emptyset$ .

## because

```
int read_len = read(fd, buf, size); // taint source
if(read_len > 0){
    printf("read data: %s\n", buf);
    for (int i = 0; i<2; i++){
        buf[i] = i;
        write(fd, buf[i], 1); // taint sink
    }
</pre>
```

0x8048541 (0x80484fd) toy\_test+0x44 8048541: mov edx, dword ptr [ebp - 0x10] **8048544**: mov eax, dword ptr [ebp + 0xc] 8048547: add edx, eax **8048549**: mov eax, dword ptr [ebp – 0x10] 804854c: mov byte ptr [edx], al **804854e**: mov edx, dword ptr [ebp – 0x10] **8048551**: mov eax, dword ptr [ebp + 0xc] **8048554**: add eax, edx 8048556: movzx eax, byte ptr [eax] 8048559: movsx eax, al **804855c**: mov dword ptr [esp + 8], 1 **8048564**: mov dword ptr [esp + 4], eax **8048568**: mov eax, dword ptr [ebp + 8] 804856b: mov dword ptr [esp], eax 804856e: call 0x80483f0 <write@plt>

basic block B (bbl-b)

The data delivered to the sink are irrelevant to the tainted data from the source.





- Tracking policy construction
  - Slow path →toy\_test(bbl-a)
    - Contain potentially tainted instructions
    - Not hot *or* TaintedMem(bbl) ∩ MergedDep(bbl) ≠ Ø

## because

0x80484fd (0x80484fd) toy_test
80484fd: push ebp
80484fe: mov ebp, esp
8048500: sub esp, 0x28
8048503: mov eax, dword ptr [ebp + 0x10]
8048506: mov dword ptr [esp + 8], eax
<b>804850a</b> : mov eax, dword ptr [ebp + 0xc]
<b>804850d</b> : mov dowrd ptr [esp + 4], eax
8048511: mov eax, dword ptr [ebp + 8]
8048514: mov dword ptr [esp], eax
<b>8048517</b> : call 0x8048390 <read@plt></read@plt>



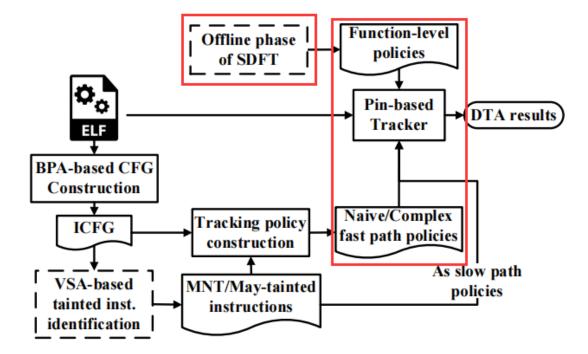
- PDG-based function abstract
  - Function fast path  $\rightarrow$  printf etc..
- Pin-based Tracker

#### Input:

- Function-level policies
- Naive/Complex fast path policies
- Slow path policies

#### output:

• DTA results



#### Next — podft's efficiency and effectiveness.



**Evaluations** 

- Experimental Settings
  - Desktop with a 2.8GHz×4 Intel Core(TM) i7-7700HQ CPU, 8GB RAM, and Linux 3.16.0 kernel (Ubuntu 14.04 32-bit).
  - The DBI framework is Pin v2.14, and libdft.

• Benchmark Programs

TABLE II. BENCHMARK PROGRAMS

Dataset ID	Description
S1	10 SPEC CPU 2k6 benchmarks
S2	3 server programs, i.e., Nginx (1.22.0), Apache httpd (2.4.7),
	and MySQL (5.5.62)
<b>S</b> 3	9 CVE programs, as presented in Table V



#### Efficiency of podft

• Compare podft's efficiency with Taint Rabbit, Dytan, Triton, and Taintgrind.

## podft achieves slowdowns of 1.6x to 27.9x with an average slowdown of 10.6x. podft is more

efficient than the

other DTA tools.

TABLE III.PERFORMANCE IMPROVEMENT OF PODFT COMPARED WITH OTHER DTA APPROACHES ON DATASET \$1 AND \$2. SLOWDOWN<br/> $SD_{TOOL} = T_{TOOL}/T_{ORIG}$ , S.T. TOOL=PODFT, TAINTRABBIT-ID, TAINTRABBIT-BV, TAINTGRIND, DYTAN, TRITON

	$T_{\rm orig}$	ig podft		TaintRabbit-ID		TaintRabbit-BV		Taintgrind		Dytan		Triton	
Benchmark	(s)	$T_{podft}(s)$	SDpodft	$T_{\rm ID}(s)$	SD <sub>ID</sub>	$T_{\rm BV}({\rm s})$	SD <sub>BV</sub>	$T_{\text{Taintgrind}}(s)$	SD <sub>Taintgrind</sub>	T <sub>Dytan</sub> (s)	SD <sub>Dytan</sub>	$T_{\text{Triton}}(s)$	SD <sub>Triton</sub>
400.perlbench	4.3	120.1	27.9	126.6	29.4	204.9	47.7	346.7	80.6	3,208.6	746.2	Timeout	N/A
401.bzip2	5.4	46.4	8.6	85.7	15.9	91.3	16.9	1,634.9	302.8	20,812.6	3,854.2	Timeout	N/A
429.mcf	2.5	5.2	2.1	34.2	13.7	Timeout	N/A	193.6	77.4	2,845.4	1,138.2	Timeout	N/A
445.gobmk	15.3	128.2	8.4	219.2	14.3	204.7	13.4	6,039.4	394.7	45,147.5	2,950.8	Timeout	N/A
456.hmmer	2.2	55.6	25.3	74.8	34.0	78.5	35.7	998.6	453.9	11,531.6	5,241.6	Timeout	N/A
458.sjeng	4.6	28.3	6.2	34.9	7.6	39.2	8.5	1,355.8	294.7	Failed	N/A	Timeout	N/A
462.libquantum	0.5	3.6	7.2	6.8	13.6	8.2	16.4	22.9	45.8	333.2	666.4	Timeout	N/A
464.h264ref	12.6	191.8	15.2	762.3	60.5	2,588.7	205.5	5,057.7	401.4	81,798.6	6,492.0	Timeout	N/A
471.omnetpp	0.4	9.1	22.8	54.6	136.5	54.1	135.3	259.6	649.0	1,401.5	3,503.8	Timeout	N/A
473.astar	9.1	14.5	1.6	129.6	14.2	129.2	14.2	1,467.4	161.3	19,525.1	2,145.6	Timeout	N/A
mysqlslap_myisam	0.6	4.5	7.5	27.5	45.8	404.2	673.7	77.5	129.2	951.9	1,586.5	Timeout	N/A
mysqlslap_innodb	0.7	4.9	7.0	28.8	41.1	392.3	560.4	73.8	105.4	999.1	1,427.3	Timeout	N/A
httpd_req_10k	0.6	5.3	8.8	11.7	19.5	15.9	26.5	26.7	44.5	227.6	379.3	40,509.2	67,515.3
httpd_req_100k	4.9	27.4	5.6	39.8	8.1	74.4	15.2	219.3	44.8	2,386.9	487.1	Timeout	N/A
Nginx_req_10k	0.5	5.5	11.0	9.8	19.6	11.7	23.4	23.9	47.8	201.1	402.2	33,671.6	67,343.2
Nginx_req_100k	4.2	20.1	4.8	24.7	5.9	32.8	7.8	196.4	46.8	2,062.3	491.0	Timeout	N/A
Avg. slowdown		-	10.6	-	30.0	-	120.0	-	205.0	-	2,100.8	-	67,429.3



#### Efficiency of podft

• Compare podft's efficiency with SELECTIVETAINT.

#### TABLE IV. PERFORMANCE COMPARISON OF PODFT WITH SELECTIVETAINT

	podft	STATICT	AINTALL	SELECTIVETAINT		
Benchmark	SD <sub>podft</sub>	T <sub>All</sub> (s)	SD <sub>All</sub>	$T_{\text{Select}}(s)$	SD <sub>Select</sub>	
400.perlbench	27.9	26.8	6.2	24.2	5.6	
401.bzip2	8.6	263.1	48.7	174.1	32.2	
429.mcf	2.1	39.1	15.6	3.8	1.5	
445.gobmk	8.4	543.4	35.5	538.5	35.2	
456.hmmer	25.3	161.4	73.4	129.2	58.7	
458.sjeng	6.2	126.8	27.6	94.8	20.6	
462.libquantum	7.2	6.3	12.6	4.5	9.0	
464.h264ref	15.2	1,255.9	99.7	801.2	63.6	
471.omnetpp	22.8	16.4	41.0	14.9	37.3	
473.astar	1.6	277.4	30.5	61.7	6.8	
Avg. slowdown	12.5	-	39.1	-	27.1	

**podft** achieves slowdowns of 1.6x to 27.9x with an average slowdown of 12.5x. and is generally more efficient than SELECTIVETAINT.



### • Effectiveness of podft's Dynamic Taint Analysis

## Real exploits detection by podft

Develop Pintool over podft to track vulnerability of CVEs

TABLE V.	EFFECT OF PODFT ON TRACKING VULNERABILITIES OF
CVEs (RCE	E=REMOTE CODE EXECUTION, S-OF=STACK OVERFLOW,
	H-OF=HEAP OVERFLOW)

ID	Program	Туре	$T_{podft}(\mathbf{s})$
CVE-2021-41253	Zydis v3.2.0	H-OF	1.4
CVE-2019-8354	SoX v14.4.2	H-OF	2.8
CVE-2018-19655	dcraw v9.28	S-OF	1.7
CVE-2018-11575	ngiflib v0.4	S-OF	1.4
CVE-2018-6612	jhead v3.00	H-OF	1.2
CVE-2017-1000437	Gravity v0.3.5	RCE	9.8
CVE-2017-14411	MP3Gain v1.5.2	S-OF	2.2
CVE-2013-2028	Nginx v1.4.0	S-OF	1.1



Work in progress

### more scalable and more flexible to be used in traditional DTA

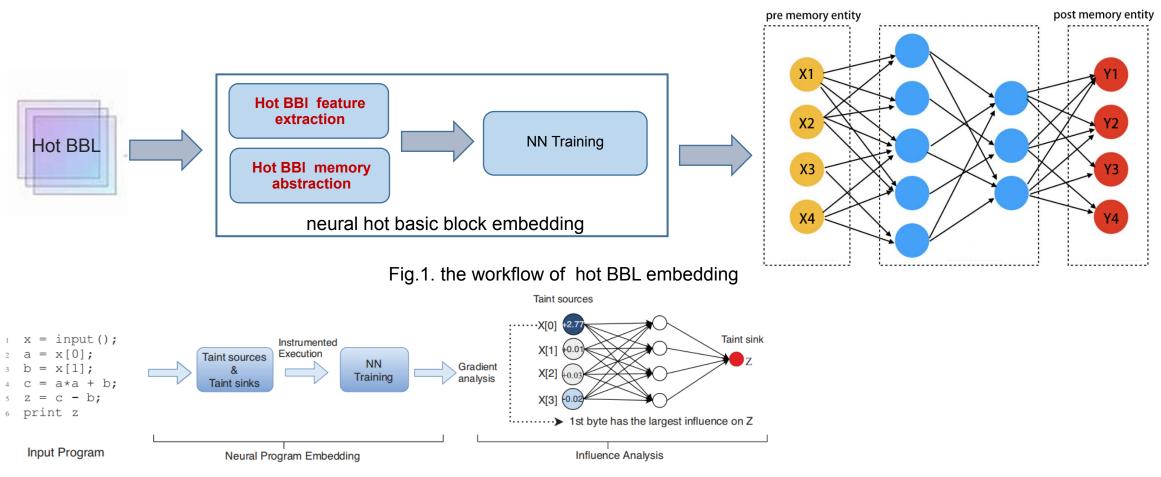


Fig.2. the workflow of NeuTaint(SP2020)





#### **Thanks for listening**